8086

16-Bit Microprocessor iAPX86 Family FINAL

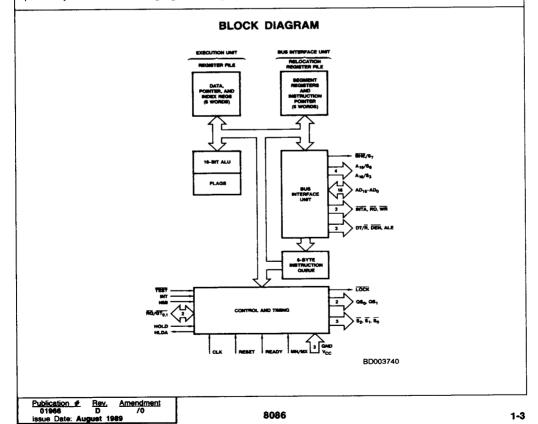
DISTINCTIVE CHARACTERISTICS

- Directly addresses up to 1 Mbyte of memory
- 24 operand addressing modes
- Efficient implementation of high level languages
- Instruction set compatible with 8080 software
- · Bit, byte, word, and block operations
- 8 and 16-bit signed and unsigned arithmetic in binary or decimal
- MULTIBUS® system interface
- Three speed options
 - 5MHz for 8086
 - 8MHz for 8086-210MHz for 8086-1

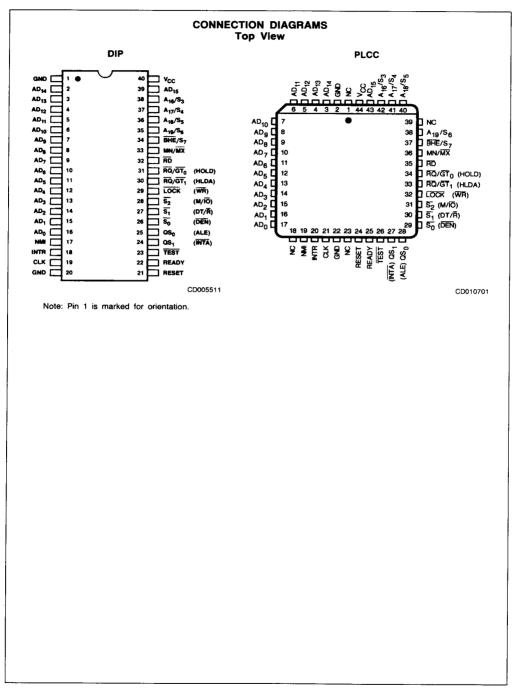
GENERAL DESCRIPTION

The 8086 is a general purpose 16-bit microprocessor CPU. Its architecture is built around thirteen 16-bit registers and nine 1-bit flags. The CPU operates on 16-bit address spaces and can directly address up to 1 megabyte using offset addresses within four distinct memory segments, designated as code, data, stack and extra code. The 8086 implements a powerful instruction set with 24 operand addressing modes. This instruction set is compatible with that of the 8080 and 8085. In addition, the 8086 is particularly effective in executing high level languages.

The 8086 can operate in minimum and maximum modes. Maximum mode offloads certain bus control functions to a peripheral device and allows the CPU to operate efficiently in a multi-processor system. The CPU and its high performance peripherals are MULTIBUS compatible. The 8086 is implemented in N-channel, depletion load, silicon gate technology and is contained in a 40-pin CERDIP package, Molded DIP package, or Plastic Leaded Chip Carrier.





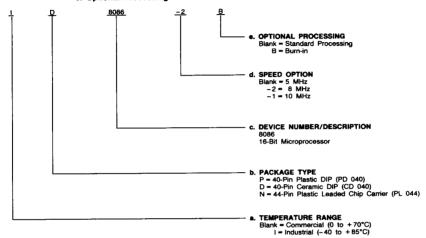


ORDERING INFORMATION

Commercial Products

AMD commercial products are available in several packages and operating ranges. The order number (Valid Combination) is formed by a combination of: a. Temperature Range

- b. Package Type
- c. Device Number
- d. Speed Option
- e. Optional Processing



Valid Combinations					
	8086				
P, D, N	8086-2				
	8086-1				
D, ID	8086-2B				
D	8086-1				
ID	8086B				

Valid Combinations

Valid Combinations list configurations planned to be supported in volume for this device. Consult the local AMD sales office to confirm availability of specific valid combinations, to check on newly released valid combinations, and to obtain additional data on AMD's standard military grade products.

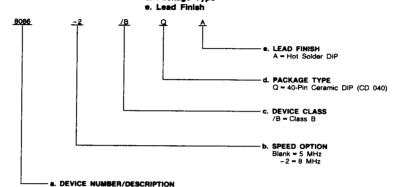


ORDERING INFORMATION

Military Products

AMD products for Aerospace and Defense applications are available in several packages and operating ranges. APL (Approved Products List) products are fully compliant with MIL-STD-883C requirements. The order number (Valid Combination) for APL products is formed by a combination of: a. Device Number

- b. Speed Option (if applicable)
- c. Device Class
- d. Package Type



| Valid Combinations | 8086 | 8086-2 | /BQA |

16-Bit Microprocessor iAPX Family

Valid Combinations

Valid Combinations list configurations planned to be supported in volume for this device. Consult the local AMD sales office to confirm availability of specific valid combinations or to check for newly released valid combinations.

Group A Tests

Group A tests consist of Subgroups 1, 2, 3, 7, 8, 9, 10, 11.

PIN DESCRIPTION

The following pin function descriptions are for 8086 systems in either minimum or maximum mode. The "Local Bus" in these descriptions is the direct multiplexed bus interface connection to the 8086 (without regard to additional bus buffers).

Pin No.*	Name	1/0	Description							
39, 2-16	AD ₁₅ -AD ₀	1/0	Address Data Bus. These lines constitute the time multiplexed memory/IO address (T ₁) and data (T ₂ , T ₃ , T _w , T ₄ , bus. A ₀ is analogous to BHE for the lower byte of the data bus, pine D ₂ -D ₀ . It is LOW during T ₁ when a byte is to be transferred on the lower portion of the bus in memory or I/O operations. Eight-bit oriented devices tied to the lower half would normally use A ₀ to condition chip select functions. (See BHE, These lines are active HIGH and float to three-state OFF during interrupt acknowledge and local bus "Hold acknowledge."							
35–38	A19/Se, A18/S5, A17/S4, A16/S3	0	Address/Status. During T_1 these are the four most significant address lines for memory operations. During I/C operations these lines are LOW. During memory and I/O operations, status information is available on these lines during T_2 , T_3 , T_4 , and T_4 . The status of the interrupt enable FLAG bit (S_2) is updated at the beginning of each CLF cycle. T_4 and T_4 , T_4 are encoded as shown. This information indicates which relocation register is presently being used for data accessing. These lines float to three-state OFF during local bus "hold acknowledge."							
		1	A ₁₇ /S ₄ A ₁₈ -S ₃ Characteristics							
			0 (LOW) 0 Alternate Data							
		1	0 1 Stack							
			1 (HIGH) 0 Code or None							
		1	1 1 Data							
			Se is 0 (LOW)							
34	BHE/S ₇	0	is High Enable/Status. During T ₁ the bus high enable signal (BHE) should be used to enable data onto e most significant half of the data bus, pins D ₁₅ -D ₈ . Eight-bit oriented devices teed to the upper half of e bus would normally use BHE to condition chip select functions. BHE is LOW during T ₁ for read, write di interrupt acknowledge cycles when a byte is to be transferred on the high portion of the bus. The Status information is available during T ₂ , T ₃ , and T ₄ . The signal is active LOW and floats to resestate OFF in "hold." It is LOW during T ₁ for the first interrupt acknowledge cycle.							
		1	BHE A ₀ Characteristics							
		1	0 0 Whole word							
			0 1 Upper byte from/ to odd address							
			1 0 Lower byte from/ to even address							
			1 1 None							
32	FID	0	Read. Read strobe indicates that the processor is performing a memory of I/O read cycle, depending on the state of the S2 pin. This signal is used to read devices which reside on the 8088 local bus. RD is active LOW during T2. T3, and Tw of any read cycle and is guaranteed to remain HIGH in T2 until the 8088 local bus has floated. This signal floats to three-state OFF in "hold acknowledge."							
22	READY		READY. Is the acknowledgment from the addressed memory or I/O device that it will complete the data transfer. The READY signal from memoryI/O is syncthronized by the 8284A Clock Generator to form READY. This signal is active HIGH. The 8086 READY input is not synchronized. Correct operation is not guaranteed if the set-up and hold times are not met.							
18	INTR	1	Interrupt Request. Is a level triggered input which is sampled during the last clock cycle of each instruc- tion to determine if the processor should enter into an interrupt acknowledge operation. A subroutine is vectored to via an interrupt vector lookup table located in system memory, it can be internally masked by software resetting the interrupt enable bit. INTR is internally synchronized. This signal is active HIGH.							
23	TEST	'	software resetting the intertion enables it. Instruction. If the TEST input is LOW, execution continues; other- wise, the processor waits in an "Idle" state. This input is synchronized internally during each clock cycle on the leading edge of CLK.							
17	NMI	'	Non-Maskable Interrupt. An edge-triggered input which causes a type 2 interrupt. A subroutine is vectored to via an interrupt vector lookup table located in system memory. NMI is not maskable internally by software. A transition from a LOW to HIGH initiates the interrupt at the end of the current instruction. Th input is internally synchronized.							
21	RESET	'	Reset. Causes the processor to immediately terminate its present activity. The signal must be active HIGI for at least four clock cycles. It restarts execution, as described in the Instruction Set description, when RESET is internally synchronized.							
19	CLK	1	Clock. Provides the basic timing for the processor and bus controller. It is asymmetric with a 33% duty cycle to provide optimized internal timing.							
40	Vcc		V _{CC} . The + 5 V power supply pin.							
1, 20	GND		Ground. The ground pin.							
	MN/MX		Minimum/Maximum. Indicates what mode the processor is to operate in. The two modes are discussed in							

8086



31, 30 RO/GTo, I/O Request at the epriority to grant see Request state Reduction Reducti	Islae of 1 CLK wide from another local bus master indicates a local bus request ("hold") to the 3 (pulse 1). 9 a T4 or T1 clock cycle, a pulse 1 CLK wide from the 8086 to the requesting master (pulse 2), ates that the 8086 has allowed the local bus to float and that it will enter the "hold acknowledge" at the next CLK. The CPU's bus interface unit is disconnected logically from the local bus during d acknowledge." Julse 1 CLK wide from the requesting master indicates to the 8086 (pulse 3) that the "hold" est is about to end and that the 8086 can reclaim the local bus at the next CLK. Taster-master exchange of the local bus is a sequence of 3 pulses. There must be one dead CLK flater each bus exchange. Pulses are active LOW.						
31, 30	\$\frac{\frac						
31, 30 RO/GTo, RG/GTo, RG/GTTo I/O Request at the e priority t grant se at the e priority t grant se state "hold 3. A pul reque Each mr cycle aff if the re of the c 1. Request 2. Currer 3. Currer 4. A local If the local 2. A me with c 2. Currer 4. A local If the local 2. A me with c 2. Currer 5. Currer 6. A local 2. A me with c 2. Currer 6. A local 2. A me and the local 3. A me and the local 3	(LOW) 0 0 Interrupt Acknowledge Acknowledge 0 1 Read I/O Port 1 0 Write I/O Port 1 1 1 Halt (HIGH) 0 0 Code Access 0 1 1 Read Memory 1 1 1 Passive						
31, 30 RO/GTo. I/O Request at the eprionity it grant se state 1. A pull 8086 2. Durin, indica state 1. Appl 1. A	0 1 Read I/O Port 1 0 Write I/O Port 1 1 1 Hait (HIGH) 0 0 Code Access 0 1 1 Read Memory 1 1 0 Write Memory 1 1 0 Write Memory 1 1 1 1 Passive st/Grant. Pins are used by other local bus masters to force the processor to release the local bus end of the processor's current bus cycle. Each pin is bidirectional with RG/GT ₀ having higher than RG/GT ₁ . RG/GT has an internal pull-up resistor so it may be left unconnected. The request/equence is as follows: 13se of 1 CLK wide from another local bus master indicates a local bus request ("hold") to the 3 (pulse 1). 19 a T ₄ or T ₁ clock cycle, a pulse 1 CLK wide from the 8086 to the requesting master (pulse 2), ates that the 8086 has allowed the local bus to float and that it will enter the "hold acknowledge" at the next CLK. The CPU's bus interface unit is disconnected logically from the local bus during dischowledge." Use 1 CLK wide from the requesting master indicates to the 8086 (pulse 3) that the "hold" est is about to end and that the 8086 can reclaim the local bus at the next CLK. Insert each bus exchange. Pulses are active LOW.						
31, 30 RO/GTo. I/O Request at the eprionity it grant se state 1. A pull 8086 2. Durin, indica state 1. Appl 1. A	1 0 Write I/O Port 1 1 1 Halt (HIGH) 0 0 Code Access 0 1 Read Memory 1 1 0 Write Memory 1 1 1 Passive st/Grant. Pins are used by other local bus masters to force the processor to release the local bus end of the processor's current bus cycle. Each pin is bidirectional with RO/GTo having higher than RO/GTo, having nigher than RO/GTO, having nigher than RO/GTO, having nigher request ("hold") to the stopping of the processor's current bus cycle. Each pin is bidirectional with RO/GTO, having higher than RO/GTO, having nigher request ("hold") to the stopping of the processor's current bus cycle. Each pin is bidirectional with RO/GTO, having higher than RO/GTO, having higher request ("hold") to the stopping of the local bus to float and that it will enter the "hold acknowledge" at the next CLK. The CPU's bus interface unit is disconnected logically from the local bus during dacknowledge." Julies 1 CLK wide from the requesting master indicates to the 8086 (pulse 3) that the "hold" est is about to end and that the 8086 can reclaim the local bus at the next CLK. Inaster-master exchange of the local bus is a sequence of 3 pulses. There must be one dead CLK fleer each bus exchange. Pulses are active LOW.						
31, 30 RG/GTo, RG/GTTo, RG	(HIGH) 0 0 Code Access 0 1 Read Memory 1 0 Write Memory 1 1 1 Passive st/Grant. Pins are used by other local bus masters to force the processor to release the local bus end of the processor's current bus cycle. Each pin is bidirectional with RO/GT ₀ having higher than RO/GT ₀ . RO/GT has an internal pull-up resistor so it may be left unconnected. The request/equence is as follows: ulse of 1 CLK wide from another local bus master indicates a local bus request ("hold") to the 3 pulse 1 CLK wide from the 8086 to the requesting master (pulse 2). ng a T ₄ or T ₁ clock cycle, a pulse 1 CLK wide from the 8086 to the requesting master (pulse 2), attes that the 8086 has allowed the local bus to float and that it will enter the "hold acknowledge" at the next CLK. The CPU's bus interface unit is disconnected logically from the local bus during dacknowledge." ulse 1 CLK wide from the requesting master indicates to the 8086 (pulse 3) that the "hold" est is about to end and that the 8086 can reclaim the local bus at the next CLK. The caster-master exchange of the local bus is a sequence of 3 pulses. There must be one dead CLK filter each bus exchange. Pulses are active LOW.						
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1 1 1 1 1 1 1 1 1 1	1 1 Passive st/Grant. Pins are used by other local bus masters to force the processor to release the local bus end of the processor's current bus cycle. Each pin is bidirectional with RO/GT ₀ having higher than RO/GT ₁ . RO/GT has an internal pull-up resistor so it may be left unconnected. The request equence is as follows: ulse of 1 CLK wide from another local bus master indicates a local bus request ("hold") to the 5 (pulse 1). ng a T ₄ or T ₁ clock cycle, a pulse 1 CLK wide from the 8086 to the requesting master (pulse 2), ates that the 8086 has allowed the local bus to float and that it will enter the "hold acknowledge at the next CLK. The CPU's bus interface unit is disconnected logically from the local bus during d acknowledge." ulse 1 CLK wide from the requesting master indicates to the 8086 (pulse 3) that the "hold" est is about to end and that the 8086 can reclaim the local bus at the next CLK. naster-master exchange of the local bus is a sequence of 3 pulses. There must be one dead CLK fittle reach bus exchange. Pulses are active LOW.						
31, 30 RG/GTo, RG/GTTo Request at the e priority t grant se en priority t grant se en	st/Grant. Pins are used by other local bus masters to force the processor to release the local bus end of the processor's current bus cycle. Each pin is bidirectional with RO/GT ₀ having higher than RO/GT ₁ . RQ/GT has an internal pull-up resistor so it may be left unconnected. The request equence is as follows: ulse of 1 CLK wide from another local bus master indicates a local bus request ("hold") to the 5 (pulse 1). ng a T ₄ or T ₁ clock cycle, a pulse 1 CLK wide from the 8086 to the requesting master (pulse 2), ates that the 8086 has allowed the local bus to float and that it will enter the "hold acknowledge at the next CLK. The CPU's bus interface unit is disconnected logically from the local bus during d acknowledge." ulse 1 CLK wide from the requesting master indicates to the 8086 (pulse 3) that the "hold" est is about to end and that the 8086 can reclaim the local bus at the next CLK. asster-master exchange of the local bus is a sequence of 3 pulses. There must be one dead CLK fler each bus exchange. Pulses are active LOW.						
## HO/GT1 at the e priority t grant se 1. A pul 8086 2. Durinn indicate state "hold 3. A pul Preque Each marcycle aff ## the re of the c 1. Reque 2. Currer 3. Currer 4. A loci If the loci 1. Local 2. A me with c 29	Islae of 1 CLK wide from another local bus master indicates a local bus request ("hold") to the 3 (pulse 1). No a T4 or T1 clock cycle, a pulse 1 CLK wide from the 8086 to the requesting master (pulse 2), ates that the 8086 has allowed the local bus to float and that it will enter the "hold acknowledge at the next CLK. The CPU's bus interface unit is disconnected logically from the local bus during d acknowledge." Use 1 CLK wide from the requesting master indicates to the 8086 (pulse 3) that the "hold" est is about to end and that the 8086 can reclaim the local bus at the next CLK. The CPU's bus is a sequence of 3 pulses. There must be one dead CLK fitter each bus exchange. Pulses are active LOW.						
Feque Each may Cycle aff If the record If the local	est is about to end and that the 8086 can reclaim the local bus at the next CLK. naster-master exchange of the local bus is a sequence of 3 pulses. There must be one dead CLK fter each bus exchange. Pulses are active LOW.						
If the re of the c	•						
1. Requerate 2. Currer 3. Currer 4. A loci 1. Local 2. A me 2. Currer 4. A loci 1. Local 2. A me	equest is made while the CPU is performing a memory cycle, it will release the local bus during T cycle when all the following conditions are met:						
1. Local 2. A me with color	uest occurs on or before T ₂ . ant cycle is not the low byte of a word (on an odd address). ent cycle is not the first acknowledge of an interrupt acknowledge sequence. cked instruction is not currently executing.						
COCK is until the "hold ac 24, 25	the local bus is idle when the request is made, two possible events will follow: Local bus will be released during the next clock. A memory cycle will start within 3 clocks. Now the four rules for a currently active memory cycle appl with condition number 1 already satisfied.						
28 M/TO O Status lir from an T4 of the	OCK. Output indicates that other system bus masters are not to gain control of the system bus while OCK is active LOW. The LOCK signal is activated by the "LOCK" prefix instruction and remains active intil the completion of the next instruction. This signal is active LOW, and floats to three-state OFF in hold acknowledge."						
28 M/IO O Status lir from an T ₄ of the							
from an T ₄ of the	QS ₁ and QS ₀ provide status to allow external tracking of the internal 8086 instruction queue. Status line. Logically equivalent to S ₂ in the maximum mode, it is used to distinguish a memory access						
29 WE I O Write In	Status line. Logically equivalent to S_2 in the maximum mode, it is used to distinguish a memory access from an I/O access. M/IO becomes valid in the T_4 preceding a bus cycle and remains valid until the final T_4 of the cycle (M = HIGH, IO = LOW). M/IO floats to three-state OFF in local bus "hold acknowledge."						
	ndicates that the processor is performing a write memory or write $1/O$ cycle, depending on the $1/O$ signal. WR is active for $1/O$, $1/O$, and $1/O$ write cycle. It is active LOW and floats to ate OFF in local bus "hold acknowledge."						
24 INTA O INTA is of each	s used as a read strobe for interrupt acknowledge cycles. It is active LOW during T_2 , T_3 , and T_W interrupt acknowlegde cycle.						
25 ALE O Address a HIGH	Latch Enable. Provided by the processor to latch the address into 8282/8283 address latch. It is pulse active during T_1 of any bus cycle. Note that ALE is never floated.						
er. It is i	ansmit/Receive. Needed in minimum system that desires to use an 8286/8287 data bus transceiv- used to control the direction of data flow through the transceiver. Logically DT/R is equivalent to be maximum mode, and its timing is the same as for M/IO. (T = HIGH, R = LOW.) This signal be three-state OFF in local bus "hold acknowledge."						
26 DEN O Data Ens transceiv	hable. Provided as an output enable for the 8266/8287 in a minimum system which uses the ver. DEN is active LOW during each memory and I/O access and for INTA cycles. For a read or rice, it is active from the middle of T2 until the middle of T4, while for a write cycle, it is active beginning of T2 until the middle of T4. DEN floats to three-state OFF in local bus "hold						

PIN DESCRIPTION (continued)								
Pin No.*	Name	1/0	Description					
31, 30	HOLD, HLDA	1/0	HOLD. Indicates that another master is requesting a local bus "hold." To be acknowledged, HOLD must be active HIGH. The processor receiving the "hold" request will issue HLDA (HIGH) as an acknowledgement in the middle of a T ₄ or T ₁ clock cycle. Simultaneous with the issuance of HLDA, the processor will float the local bus and control lines. After HOLD is detected as being LOW, the processor will LOWer HLDA, and when the processor needs to run another cycle, it will again drive the local bus and control lines. The same rules as for RO/GT apply, regarding when the local bus will be released. HOLD is not asynchroneous input. External synchronization should be provided if the system cannot otherwise quarantee the set-up time.					

^{*}Pin numbers correspond to DIPs only.

DETAILED DESCRIPTION

The 8086 CPU is internally organized into two processing units. These two units are the Bus Interface Unit (BIU) and the Execution Unit (EU). A block diagram of this organization is shown on page 1.

The BIU performs instruction fetch and queuing, operand fetch and store, address relocation, and basic bus control. The EU receives operands and instructions from the BIU and processes them on a 16-bit ALU. The EU accesses memory and peripheral devices through requests to the BIU. The BIU generates physical addresses in memory using the 4 segment registers and offset values.

The BIU and EU usually operate asynchronously. This permits the 8086 to overlap execution fetch and execution. Up to 6 instruction bytes can be queued. The instruction queue acts as a FIFO buffer for instructions, from which the EU extracts instruction bytes as required.

Memory Organization

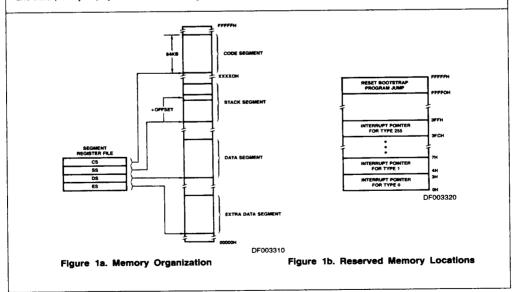
The 8086 addresses up to 1 megabyte of memory. The address space is organized as a linear array, from 00000 to FFFFF in hexadecimal. Memory is subdivided into segments of 64K bytes each. There are 4 segments: code, stack, data, and extra (usually employed as an extra data segment). Each

segment thus contains information of a similar type. Selection of a destination segment is automatically performed using the rules in the table below. This segmentation makes memory more easily relocatable and supports a more structured programming style.

Physical addresses in memory are generated by selecting the appropriate segment, obtaining the segment "base" address from the segment register, shifting the base address 4 digits to the left, and then adding this base to the "offset" address. For programming code, the offset address is obtained from the instruction pointer. For operands, the offset address is calculated in several ways, depending upon information contained in the addressing mode. Memory organization and address generation are shown in Figure 1a.

Certain memory locations are reserved for specific CPU operations. These are shown in Figure 1b. Addresses FFFFOH through FFFFFH are reserved for operations which include a jump to the initial program loading routine. After RESET, the CPU will always begin execution at location FFFFOH, where the jump must be located.

Addresses 00000H through 003FFH are reserved for interrupt operations. The service routine of each of the 256 possible interrupt types is signaled by a 4-byte pointer. The pointer elements must be stored in reserved memory addresses before the interrupts are invoked.





Memory Reference Need	Segment Register Used	Segment Selection Rule
Instructions	CODE (CS)	Automatic for all prefetching of instructions.
Stack	STACK (SS)	All stack pushes and pops, and all memory references relative to BP base register except data references.
Local Data	DATA (DS)	Data references which are relative to the stack, the destination of a string operation, or explicitly overriden.
External (Giobal) Data	EXTRA (ES)	Destination of string operations, when they are explicitly selected using a segment override.

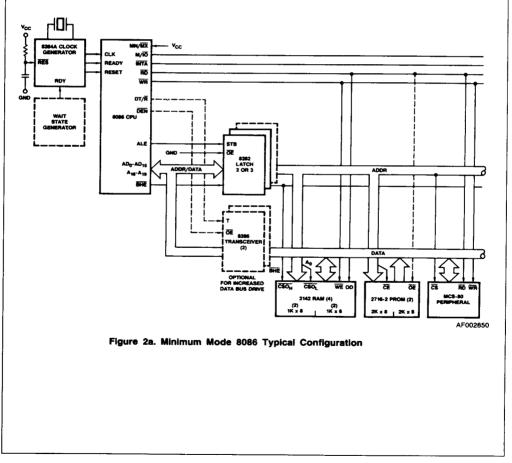
Minimum and Maximum Modes

The 8086 has two system configurations, minimum and maximum mode. The CPU has a strap pin, MN/MX, which defines the system configuration. The status of this strap pin defines the function of pin numbers 24 through 31.

When MN/MX is strapped to GND, the 8086 operates in maximum mode. The operations of pins 24 through 31 are redefined. In maximum mode, several bus timing and control functions are "off-loaded" to the 8288 bus controller, thus

freeing up the CPU. The CPU communicates status information to the 8288 through pins $S_0,\,S_1,\,$ and $S_2.\,$ In maximum mode, the 8086 can operate in a multiprocessor system, using the LOCK signal within a Multibus format.

When MN/ $\overline{\text{MX}}$ is strapped to V_{CC}, the 8086 operates in minimum mode. The CPU sends bus control signals itself through pins 24 through 31. This is shown in the Connection Diagrams (in parentheses). Examples of minimum and maximum mode systems are shown in Figure 2.



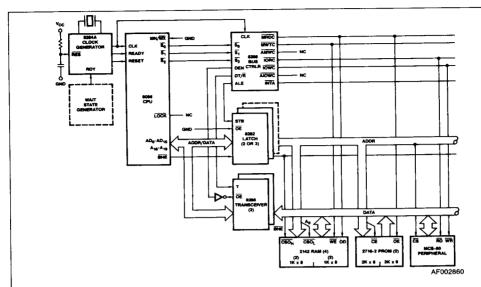


Figure 2b. Maximum Mode 8086 Typical Configuration

Bus Operation

The 8086 has a combined address and data bus, commonly referred to as "a time multiplexed bus." This technique provides the most efficient use of pins on the processor while permitting the use of a standard 40-lead package. This bus can be used throughout the system with address latching provided on memory and I/O modules. The bus can also be demultiplexed at the processor with a single set of address latches if a standard non-multiplexed bus is desired for the system.

Each bus cycle consists of at least four CLK cycles. These are referred to as T_1 , T_2 , T_3 and T_4 (see Figure 5). The address is sent from the processor during T_1 . Data transfer occurs on the bus during T_3 and T_4 . T_2 is used for changing the direction of the bus during read operations. In the event that a "NOT READY" indication is given by the addressed device, "Wait" states (T_W) are inserted between T_3 and T_4 . Each inserted "Wait" state is of the same duration as a CLK cycle. "Idle" states (T_1) or inactive CLK cycles can occur between 8086 bus cycles. The processor uses these cycles for internal housekeeping.

During T₁ of any bus cycle, the ALE (Address Latch Enable) signal is emitted (by either the processor or the 8288 bus controller, depending on the MN/ $\overline{\text{MX}}$ strap). At the trailing edge of this pulse, a valid address and certain status information for the cycle may be latched.

Status bits $\overline{S_0}$, $\overline{S_1}$, and $\overline{S_2}$ are used, in maximum mode, by the bus controller to identify the type of bus transaction according to the following table:

Ī2	Ī5₁	Ŝ₀	Characteristics
(LOW)	0	0	Interrupt Acknowledge
0	0	1	Read I/O
0	1	0	Write I/O
0	1	1	Hait
1(HIGH)	0	0	Instruction Fetch
1	0	1	Read Data from Memory
1	1	0	Write Data to Memory
1	1	1	Passive (no bus cycle)

Status bits S_3 through S_7 are multiplexed with high-order address bits and the \overline{BHE} signal, and are therefore valid during T_2 through T_4 . S_3 and S_4 indicate which segment register (see Instruction Set description) was used for this bus cycle in forming the address, according to the following table:

S ₄	S ₃	Characteristics
0(LOW)	0	Alternate Data (extra segment)
0	1	Stack
1(HIGH)	0	Code or None
1	1	Data

 S_5 is a reflection of the PSW interrupt enable bit. S_6 = 0 and S_7 is a spare status bit.

I/O Addressing

8086 I/O operations can address up to a maximum of 64K I/O byte registers or 32K I/O word registers. The I/O address appears in the same format as the memory address on bus lines A₁₅-A₀. The address lines A₁₉-A₁₆ are zero in I/O operations. I/O instructions which use register DX as a pointer have full address capability. Direct I/O instructions directly address one or two of the 256 I/O byte locations in page 0 of the I/O address space. I/O ports are addressed in the same manner as memory locations.



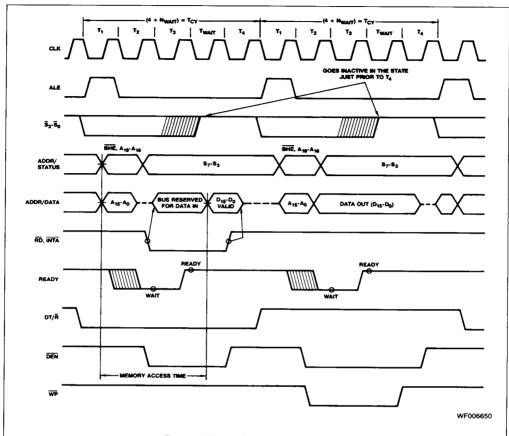


Figure 3. Basic System Timing

EXTERNAL INTERFACE

Processor Reset and Initialization

Processor initialization or start up is accomplished with activation (HIGH) of the RESET pin. The 8086 RESET is required to be HIGH for greater than 4 CLK cycles. The 8086 will terminate operations on the high-going edge of RESET and will remain dormant as long as RESET is HIGH. The low-going transition of RESET triggers an internal reset sequence for approximately 10 CLK cycles. After this interval the 8086 operates normally beginning with the instruction in absolute location FFFF0H (see Figure 1b). The details of this operation are explained in the Instruction Set description of the MCS-86 Family User's Manual. The RESET input is internally synchronized to the processor clock. At initialization the HIGH-to-LOW transition of RESET must occur no sooner than 50 μ s after power-up, to allow complete initialization of the 8086.

NMI may not be asserted prior to the 2nd CLK cycle following the end of RESET.

Interrupt Operations

Interrupt operations fall into two classes: software or hardware initiated. The software initiated interrupts and software aspects of hardware interrupts are described in the Instruction Set description. Hardware interrupts are either non-maskable or maskable.

Interrupts transfer control to a new program location. A 256-element table containing address pointers to the interrupt service program locations resides in absolute locations 0 through 3FFH (see Figure 1b), which are reserved for this purpose. Each element in the table is 4 bytes in size and corresponds to an interrupt "type." An interrupting device supplies an 8-bit type number during the interrupt acknowledge sequence, which is used to "vector" through the appropriate element to the new interrupt service program location.

Non-Maskable Interrupt (NMI)

The processor provides a single non-maskable interrupt pin (NMI) which has higher priority than the maskable interrupt request pin (INTR). A typical use would be to activate a power

failure routine. The NMI is edge-triggered on a LOW-to-HIGH transition. The activation of this pin causes a type 2 interrupt. (See Instruction Set description.)

NMI is required to have a duration in the HIGH state of greater than two CLK cycles, but is not required to be synchronized to the clock. Any high-going transition of NMI is latched on-chip and will be serviced at the end of the current instruction or between whole moves of a block-type instruction. Worst case response to NMI would be to multiply, divide, and variable shift instructions. There is no specification on the occurrence of the low-going edge; it may occur before, during, or after the servicing of NMI. Another high-going edge triggers another response if it occurs after the start of the NMI procedure. The signal must be free of logical spikes in general and be free of bounces on the low-going edge to avoid triggering extraneous responses.

Maskable Interrupt (INTR)

The 86/10 provides a single interrupt request input (INTR) which can be masked internally by software with the resetting of the interrupt enable FLAG status bit. The interrupt request signal is level-triggered. It is internally synchronized during each clock cycle on the high-going edge of CLK. To be responded to, INTR must be present (HIGH) during the clock period preceding the end of the current instruction or the end of a whole move for a block-type instruction. During the interrupt response sequence, further interrupts are disabled. The enable bit is reset as part of the response to any interrupt (INTR, NMI, software interrupt, or single-step), although the FLAGS register, which is automatically pushed onto the stack, reflects the state of the processor prior to the Interrupt. Until the old FLAGS register is restored, the enable bit will be zero unless specifically set by an instruction.

During the response sequence (Figure 4), the processor executes two successive (back-to-back) interrupt acknowledge cycles. The 8086 emits the LOCK signal from T₂ of the first bus cycle until T₂ of the second. A local bus "hold" request will not be honored until the end of the second bus cycle, in the second bus cycle, a byte is fetched from the external interrupt system (e.g., 8259A PIC) which identifies the source (type) of the interrupt. This byte is multiplied by four and used as a pointer into the interrupt vector lookup table. An INTR signal left HIGH will be continually responded to within the limitations of the enable bit and sample period. The INTERRUPT RETURN instruction includes a FLAGS pop, which returns the status of the original interrupt enable bit when it restores the FLAGS.

HALT

When a software "HALT" instruction is executed, the processor indicates that it is entering the "HALT" state in one of two ways depending upon which mode is strapped. In minimum mode, the processor issues one ALE with no qualifying bus control signals. In Maximum Mode, the processor issues appropiate HALT status on $\overline{S}_2\overline{S}_1\overline{S}_0$, and the 8288 bus controller issues one ALE. The 8086 will not leave the "HALT" state when a local bus "hold" is entered while in "HALT." In this case, the processor reissues the HALT indicator. An interrupt request or RESET will force the 8086 out of the "HALT" state.

Read/Modify/Write (Semaphore) Operation Via Lock

The LOCK status information is provided by the processor when directly consecutive bus cycles are required during the execution of an instruction. This provides the processor with the capability of performing read/modify/write operations on memory (via the Exchange Register With Memory Instruction, for example) without the possibility of another system bus

master receiving intervening memory cycles. This is useful in multiprocessor system configurations to accomplish "test and set lock" operations. The LOCK signal is activated (forced LOW) in the clock cycle following the one in which the software "LOCK" prefix instruction is decoded by the EU. It is deactivated at the end of the last bus cycle of the instruction following the "LOCK" prefix instruction. While LOCK is active, a request on a RO/GT pin will be recorded and then honored at the end of the LOCK.

External Synchronization Via Test

As an alternative to the interrupts and general I/O capabilities, the 8086 provides a single software-testable input known as the TEST signal. At any time, the program may execute a WAIT instruction. If at that time the TEST signal is inactive (HIGH), program execution becomes suspended while the processor waits for TEST to become active. It must remain active for at least 5 CLK cycles. The WAIT instruction is reexecuted repeatedly until that time. This activity does not consume bus cycles. The processor remains in an idle state while waiting. All 8086 drivers go to three-state OFF if bus "HOLD" is entered. If interrupts are enabled, they may occur while the processor is waiting. When this occurs, the processor fetches the WAIT instruction one extra time, processes the interrupt, and then re-fetches and re-executes the WAIT instruction upon returning from the interrupt.

Basic System Timing

Typical system configurations for the processor operating in minimum mode and in maximum mode are shown in Figures 2a and 2b, respectively. In minimum mode, the processor emits bus control signals in a manner similar to the 8085. In maximum mode, the processor emits coded status information which the 8288 bus controller uses to generate MULTIBUS compatible bus control signals. Figure 3 illustrates the signal timing relationships.

System Timing - Minimum System

The read cycle begins in T1 with the assertion of the Address Latch Enable (ALE) signal. The trailing (low-going) edge of this signal is used to latch the address information, which is valid on the local bus at this time, into the 8282/8283 latch. The BHE and An signals address the low, high, or both bytes. From T₁ to T₄, the M/IO signal indicates a memory or I/O operation. At T2 the address is removed from the local bus and the bus goes to a high impedance state. The read control signal is also asserted at T2. The read (RD) signal causes the addressed device to enable its data bus drivers to the local bus. Some time later valid data will be available on the bus and the addressed device will drive the READY line HIGH. When the processor returns the read signal to a HIGH level, the addressed device will again 3-state its bus drivers. If a transceiver (8286/8287) is required to buffer the 8086 local bus, signals DT/R and DEN are provided by the 8086.

A write cycle also begins with the assertion of ALE and the emission of the address. The M/IO signal is again asserted to indicate a memory or I/O write operation. In the T_2 immediately following the address emission, the processor emits the data to be written into the addressed location. This data remains valid until the middle of T_4 . During T_2 , T_3 , and T_W , the processor asserts the write control signal. The write (\overline{WR}) signal becomes active at the beginning of T_2 as opposed to the read which is delayed somewhat into T_2 to provide time for the bus to float.

The \overline{BHE} and A_0 signals are used to select the proper byte(s) of the memory/IO word to be read or written according to the following table.



BHE	Ao	Characteristics
0	0	Whole word
0	1	Upper byte from/to odd address
1	0	Lower byte from/to even address
1	1	None

I/O ports are addressed in the same manner as memory location. Even addressed bytes are transferred on the D₇–D₀ bus lines and odd addressed bytes on D₁₅–D₈.

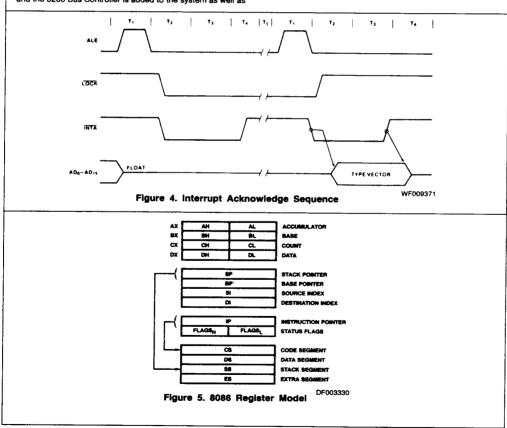
The basic difference between the interrupt acknowledge cycle and a read cycle is that the interrupt acknowledge signal (INTA) is asserted in place of the read (RD) signal and the address bus is floated. (See Figure 6.) In the second of two successive INTA cycles, a byte of information is read from bus lines D7-D0 as supplied by the interrupt system logic (i.e., 8259A Priority Interrupt Controller). This byte identifies the source (type) of the interrupt. It is multiplied by four and used as a pointer into a interrupt vector lookup table, as described earlier.

Bus Timing — Medium Size Systems

For medium size systems, the MN/ $\overline{\text{MX}}$ pin is connected to V_{SS}, and the 8288 Bus Controller is added to the system as well as

an 8282/8283 latch for latching the system address and a 8286/8287 transceiver to allow for bus loading greater than the 8086 is capable of handling. Signals ALE, DEN, and DT/R are generated by the 8288 instead of the processor in this configuration, although their timing remains relatively the same. The 8086 status (§2, §1, and §0) provide type-of-cycle information and become 8288 inputs. This bus cycle information specifies read (code, data, or I/O), write (data or I/O), interrupt acknowledge, or software halt. The 8288 thus issues control signals specifying memory read or write, I/O read or write, or interrupt acknowledge. The 8288 provides two types of write strobes, normal and advanced, to be applied as required. The normal write strobes have data valid at the leading edge of write. The advanced write strobes have the same timing as read strobes, and hence, data isn't valid at the leading edge of write. The 8286/8287 transceiver receives the usual T and OE inputs from the 8288's DT/R and DEN.

The pointer into the interrupt vector table, which is passed during the second INTA cycle, can derive from an 8259A located on either the local bus or the system bus. If the master 8259A Priority Interrupt Controller is positioned on the local bus, a TTL gate is required to disable the 8286/8287 transceiver when reading from the master 8259A during the interrupt acknowledge sequence and software "poll."



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8086

ABSOLUTE MAXIMUM RATINGS

Storage Temperature65 to +150°	С
Ambient Temperature Under Bias 0 to 70°	С
Voltage on any Pin	
with Respect to Ground1 to +7.0	٧
Power Dissipitation2.5 \	Ν

Stresses above those listed under ABSOLUTE MAXIMUM RATINGS may cause permanent device failure. Functionality at or above these limits is not implied. Exposure to absolute maximum ratings for extended periods may affect device reliability.

OPERATING RANGES

Commercial (C) Devices	
Temperature (T _A)	0 to +70°C
Supply Voltage (VCC)	
8086	5 V ± 10%
8086-1, 8086-2	5 V ± 5%
Industrial (I) Devices	
Temperature (T _A)	40 to +85°C
Supply Voltage (V _{CC})	
8086	5 V ± 10%
8086-1, 8086-2	5 V ± 5%

Operating ranges define those limits between which the functionality of the device is guaranteed.

DC CHARACTERISTICS over COMMERCIAL operating range

Parameters	Description	Test Conditions	Min	Max	Units
VIL	Input Low Voltage		-0.5	+ 0.8	٧
V _{IH}	Input High Voltage		2.0	V _{CC} + 0.5	>
VOL	Output Low Voltage	I _{OL} = 2.5 mA		0.45	>
VOH	Output High Voltage	I _{OH} = ~400 μA	2.4		>
Icc	Power Supply Current	All Speeds		340	mA
lu	Input Leakage Current	0V ≤ V _{IN} ≤ V _{CC}		±10	μА
ILO	Output Leakage Current	0.45V ≤ V _{OUT} ≤ V _{CC}		± 10	μΑ
V _{CL}	Clock Input Low Voltage		- 0.5	+ 0.6	٧
VCH	Clock Input High Voltage		3.9	V _{CC} + 1.0	V
CIN	Capacitance of Input Buffer (All Input except AD ₀ -AD ₁₅ , RQ/GT)	fc = 1 MHz		15	pF
CIO	Capacitance of I/O Buffer (AD ₀ -AD ₁₅ , RQ/GT)	fc = 1 MHz	I	15	pF

SWITCHING CHARACTERISTICS over COMMERCIAL operating range MINIMUM COMPLEXITY SYSTEM TIMING REQUIREMENTS

		Test	808	3	8086-	2	8086-	1	11-14-
Parameters	Description	Conditions	Min	Max	Min	Max	Min	Max	Units
TCLCL	CLK Cycle Period		200	500	125	500	100	500	ns
TCLCH	CLK Low Time		118		68		53		ns
TCHCL	CLK High Time		69		44		39		ns
TCH1CH2	CLK Rise Time	From 1.0 to 3.5V		10		10		10	ns
TCL2CL1	CLK Fall Time	From 3.5 to 1.0V		10		10		10	ns
TDVCL	Data in Set-up Time		30		20		5		пв
TCLDX	Data in Hold Time		10		10		10	L.,	ns
TR1VCL	RDY Set-up Time into 8284A (See Notes 1, 2)		35		35		35		ns
TCLR1X	RDY Hold Time into 8284A (See Notes 1, 2)		0		0		0	<u> </u>	ns
TRYHCH	READY Set-up Time into 8086		118		68		53	<u> </u>	ns
TCHRYX	READY Hold Time into 8086		30		20		20		ns
TRYLCL	READY Inactive to CLK (See Note 3)		-8		-8		-10		ns
THVCH	HOLD Set-up Time		35		20		20		ns
TINVCH	INTR, NMI, TEST Set-up Time (See Note 2)		30		15		15		ns
TILIH	Input Rise Time (Except CLK)	From 0.8 to 2.0 V	, , , , , , , , , , , , , , , , , , ,	20		20		20	ns
TIHIL	Input Fall Time (Except CLK)	From 2.0 to 0.8 V		12		12		12	กร

Notes: 1. Signal at 8284A shown for reference only.

2. Set-up requirement for asynchronous signal only to guarantee recognition at next CLK.

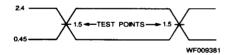
3. Applies only to T2 state (8ns into T3).



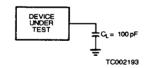
SWITCHING CHARACTERISTICS over **COMMERCIAL** operating range (continued) TIMING RESPONSES

		Test	8086		8086-2	?	8086-1	l	
Parameters	Description	Conditions	Min	Max	Min	Max	Min	Max	Units
TCLAV	Address Valid Delay		10	110	10	60	10	50	ns
TCLAX	Address Hold Time]	10		10		10		ns
TCLAZ	Address Float Delay	1	TCLAX	80	TCLAX	50	10	40	ns
TLHLL	ALE Width	}	TCLCH - 20		TCLCH-10		TCLCH - 10		ns
TCLLH	ALE Active Delay			80		50		40	ns
TCHLL	ALE Inactive Delay]		85		55		45	ns
TLLAX	Address Hold Time to ALE Inactive]	TCHCL - 10		TCHCL - 10		TCHCL - 10		ns
TCLDV	Data Valid Delay	1	10	110	10	60	10	50	กร
TCHDX	Data Hold Time	1	10		10		10		ns
TWHDX	Data Hold Time After WR]	TCLCH -30		TCLCH-30		TCLCH -25		ns
TCVCTV	Control Active Delay 1	1	10	110	10	70	10	50	ns
TCHCTV	Control Active Delay 2	*CL = 20-100 pF	10	110	10	60	10	45	ns
TCVCTX	Control Inactive Delay	for all 8086 Outputs (in addition	10	110	10	70	10	50	ns
TAZRL	Address Float to READ active	to 8086 self-load). Typical C _L = 100 pF.	0		0		0		ns
TCLRL	RD Active Delay]	10	165	10	100	10	70	ns
TCLRH	RD Inactive Delay	1	10	150	10	80	10	60	ns
TRHAV	RD Inactive to Next Address Active		TCLCL -45		TCLCL -40		TCLCL -35		ns
TCLHAV	HLDA Valid Delay]	10	160	10	100	10	60	ns
TRLRH	RD Width]	2TCLCL - 75		2TCLCL -50		2TCLCL -40		ns
TWLWH	WR Width		2TCLCL -60		2TCLCL -40		2TCLCL -35		ns
TAVAL	Address Valid to ALE Low		TCLCH - 60		TCLCH -40		TCLCH - 35		ns
TOLOH	Output Rise Time	From 0.8 to 2.0 V		20		20		20	ns
TOHOL	Output Fall Time	From 2.0 to 0.8 V		12		12		12	ns

SWITCHING TEST INPUT/OUTPUT WAVEFORM



SWITCHING TEST LOAD CIRCUIT



AC Testing inputs are driven at 2.4 V for a logic "1" and 0.45 V for a logic "0." Timing measurements are made at 1.5 V for both a logic "1" and "0."

C_L includes jig capacitance

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SWITCHING CHARACTERISTICS over COMMERCIAL operating range (continued) MAX MODE SYSTEM (USING 8288 BUS CONTROLLER) TIMING REQUIREMENTS

		Test	808	5	8086	-2	8086-	T U	11-14-
Parameters	Description	Conditions	Min	Max	Min	Max	Min	Max	Unite
TCLCL	CLK Cycle Period		200	500	125	500	100	500	ns
TCLCH	CLK Low Time		118		68		53	<u> </u>	ns
TCHCL	CLK High Time		69		44		39		ns
TCH1CH2	CLK Rise Time	From 1.0 to 3.5 V		10		10		10	ns
TCL2CL1	CLK Fall Time	From 3.5 to 1.0 V		10		10		10	ns
TDVCL	Data in Set-up Time		30		20		5		ns
TCLDX	Data in Hold Time		10		10		10		ns
TR1VCL	RDY Set-up Time into 8284A (See Notes 1, 2)		35		35		35		ns
TCLR1X	RDY Hold Time into 8284A (See Notes 1, 2)		0		0		0		ns
TRYHCH	READY Set-up Time into 8086		118		68		53		ns
TCHRYX	READY Hold Time into 8086		30		20		20		ns
TRYLCL	READY Inactive to CLK (See Note 4)		-8		-8		-10	<u> </u>	ns
TINVCH	Set-up Time for Recognition (INTR, NMI, TEST (See Note 2)		30		15		15		ns
TGVCH	RQ/GT Set-up Time		30		15		12		ns
TCHGX	RQ Hold Time into 8066		40		30		20		ns
TILIH	Input Rise Time (Except CLK)	From 0.8 to 2.0 V		20		20		20	ns
TIHIL	Input Fall Time (Except CLK)	From 2.0 to 0.8 V		12		12		12	ns

Notes: 1. Signal at 8284A or 8288 shown for reference only.

^{2.} Set-up requirement for asynchronous signal only to guarantee recognition at next CLK.

Applies only to T3 and wait states.
 Applies only to T2 state (8ns into T3).



SWITCHING CHARACTERISTICS over COMMERCIAL and INDUSTRIAL ranges (continued) TIMING RESPONSES

		Test	8086		8086-2		8086-1		l
Parameters	Description	Conditions	Min	Max	Min	Max	Min	Max	Units
TCLML	Command Active Delay (See Note 1)		10	35	10	35	10	35	ns
TCLMH	Command Inactive Delay (See Note 1)]	10	35	10	35	10	35	ns
TRYHSH	READY Active to Status Passive (See Note 3)			110		65		45	ns
TCHSV	Status Active Delay		10	110	10	60	10	45	กร
TCLSH	Status Inactive Delay	7	10	130	10	70	10	55	ns
TCLAV	Address Valid Delay		10	110	10	60	10	50	ns
TCLAX	Address Hold Time	1	10		10		10		ns
TCLAZ	Address Float Delay	7	TCLAX	80	TCLAX	50	10	40	ns
TSVLH	Status Valid to ALE High (See Note 1)			15		15		15	ns
TSVMCH	Status Valid to MCE High (See Note 1)			15		15		15	ns
TCLLH	CLK Low to ALE Valid (See Note 1)	1		15		15		15	ns
TCLMCH	CLK Low to MCE High (See Note 1)	1		15		15		15	ns
TCHLL	ALE Inactive Delay (See Note 1)	C _L = 20-100 pF for all 8086		15		15		15	ns
TCLMCL	MCE Inactive Delay (See Note 1)	Outputs (In addition to 8086 self-load)		15		15		15	ns
TCLDV	Data Valid Delay	7	10	110	10	60	10	50	ns
TCHDX	Data Hold Time	1	10		10		10		ns
TCVNV	Control Active Delay (See Note 1)	7	5	45	5	45	5	45	ns
TCVNX	Control Inactive Delay (See Note 1)	1	10	45	10	45	10	45	пѕ
TAZRL	Address Float to Read Active]	0		0		0		ns
TCLRL	RD Active Delay	1	10	165	10	100	10	70	ns
TCLRH	RD Inactive Delay	1	10	150	10	80	10	60	ns
TRHAV	RD Inactive to Next Address Active	1	TCLCL -45		TCLCL -40		TCLCL - 35		ns
TCHDTL	Direction Control Active Delay (See Note 1)			50		50		50	ns
TCHDTH	Direction Control Inactive Delay (See Note 1)			30		30		30	ns
TCLGL	GT Active Delay	7	0	85	0	50	0	38	ns
TCLGH	GT Inactive Delay	1	0	85	0	50	0	45	ns
TRLRH	RD Width	7	2TCLCL ~ 75		2TCLCL -50		2TCLCL - 40		ns
TOLOH	Output Rise Time	From 0.8 to 2.0 V		20		20		20	ns
TOHOL	Output Fall Time	From 2.0 to 0.8 V		12		12		12	ns

Notes: 1. Signal at 8284A or 8288 shown for reference only.
2. Set-up requirement for asynchronous signal only to guarantee recognition at next CLK.

^{3.} Applies only to T3 and wait states.

^{4.} Applies only to T2 state (8ns into T3).

ABSOLUTE MAXIMUM RATINGS

Storage Temperature-65 to +150°C Ambient Temperature Under Bias 0 to 70°C Voltage on any Pin with Respect to Ground-1 to +7.0 V Power Dissipitation2.5 W

Stresses above those listed under ABSOLUTE MAXIMUM RATINGS may cause permanent device failure. Functionality at or above these limits is not implied. Exposure to absolute maximum ratings for extended periods may affect device reliability.

OPERATING RANGES

Military (M) Devices	
Temperature (T _C)	55 to +125°C
Supply Voltage (VCC)	5 V ±10%

Operating ranges define those limits between which the functionality of the device is guaranteed.

DC CHARACTERISTICS over MILITARY operating range (for APL Products, Group A, Subgroups 1, 2, 3 are tested unless otherwise noted)

Parameter Symbol	Parameter Description	Test Conditions	Min.	Max.	Unit
V _{IL} †	Input LOW Voltage	V _{CC} = Min. & Max.	-0.5*	+0.8	٧
ViH †	Input HIGH Voltage	V _{CC} = Min. & Max.	2.0	V _{CC} + 0.5*	٧
VOL	Output LOW Voltage	iOL = 2.0 mA, VCC = Min.		0.45	٧
Voн	Output HIGH Voltage	l _{OH} = -400 μA, V _{CC} = Min.	2.4		٧
loc	Power Supply Current (Note 1)	T _C = 25°C, V _{CC} = Max.		340	mA
l _{LI}	Input Leakage Current	V _{CC} = Max., V _{IN} = 5.5 V & 0 V	-10	10	μΑ
ILO TT	Output Leakage Current	V _{CC} = Max., V _{OUT} = 5.5 V & 0.45 V	-10	10	μΑ
VCL †	Clock input LOW Voltage	V _{CC} = Min. & Max.	-0.5*	+ 0.6	V
Vch †	Clock input HIGH Voltage	V _{CC} = Min. & Max.	3.9	V _{CC} + 1.0*	٧
C _{IN} †††	Capacitance of Input Buffer (All Input Except ADo-AD15, RQ/GT)	fc = 1 MHz		20*	ρF
Cio ttt	Capacitance of I/O Buffer (AD ₀ -AD ₁₅ , RQ/GT)	fc = 1 MHz		20*	рF

Notes: 1. ICC is measured while running a functional pattern with spec value IOL/IOH loads applied.

Guaranteed by design; not tested.
 † Group A, Subgroups 7 and 8 only are tested.
 †† Group A, Subgroups 1 and 2 only are tested.
 ††† Not included in Group A tests.



SWITCHING CHARACTERISTICS over MILITARY operating range (for APL Products, Group A, Subgroups 9, 10, 11 are tested unless otherwise noted)

MINIMUM COMPLEXITY SYSTEM TIMING REQUIREMENTS

Parameter	Parameter	Test Conditions	80	86	808	36-2	i
Symbol	Description	(Note 6)	Min.	Max.	Min.	Max.	Unit
TCLCL	CLK Cycle Period (Note 11)		200	500	125	500	ns
TCLCH	CLK LOW Time		118		68		ns
TCHCL	CLK HIGH Time		69		44		ns
TCH1CH2	CLK Rise Time (Note 5)	From 1.0 to 3.5 V		10		10	ns
TCL2CL1	CLK Fall Time (Note 5)	From 3.5 to 1.0 V		10		10	ns
TDVCL	Data in Setup Time		30		20		ns
TCLDX	Data in Hold Time		10		10		ns
TR1VCL	RDY Setup Time into 8284A (Notes 1 & 2)		35		35		ns
TCLR1X	RDY Hold Time into 8284A (Notes 1 & 2)		0		0		ns
TRYHCH	READY Setup Time into 8086		118		68		ns
TCHRYX	READY Hold Time into 8086		30		20		ns
TRYLCL	READY Inactive to CLK (Note 4)		-8		-8		ns
THVCH	HOLD Setup Time		35		20		ns
TINVCH	INTR, NMI, TEST Setup Time (Note 2)		30		15		ns
TILIH	Input Rise Time (Except CLK) (Note 5)	From 0.8 to 2.0 V		20		20	ns
TIHIL	Input Fall Time (Except CLK) (Note 5)	From 2.0 to 0.8 V		12		12	ns

Notes:

1. Signal at 8284A and 8288 shown for reference only.

2. Setup requirement for asynchronous signal only to guarantee recognition at next CLK.

3. Applies only to T3 and wait states.

4. Applies only to T2 state (8 ns into T3).

5. Not tested: these specs are controlled by the Teradyne J941 tester.

6. Voc. = 4.5 V, 5.5 V VIH = 2.4 V
VIL = .45 V VIHC = 4.3 V
VILC = .25 V VOH = 1.6 V

7. Minimum spec tested at Voc. Max. (5.5 V) only.

8. Maximum spec tested at Voc. Min. (4.5 V) only.

10. Tested at Voc. Min. (4.5 V) only.

11. Test conditions for TCLCL Max. are:

Voc. = 4.5 V VIHC = 1 V
VIL = 0 V VIHC = 5 V



SWITCHING CHARACTERISTICS over MILITARY operating range (continued) TIMING RESPONSES

		Test	80	86	808	6-2	
Parameter Symbol	Parameter Description	Conditions (Note 6)	Min.	Max.	Min.	Max.	Unit
TCLAV	Address Valid Delay		10	110	10	60	пѕ
TCLAX	Address Hold Time (Notes 7 & 8)	ĺ	10		10		ns
TCLAZ	Address Float Delay (Note 8)		10	80	10	50	ns
TLHLL	ALE Width (Note 10)		98		58		ns
TCLLH	ALE Active Delay (Note 8)			80		50	ns
TCHLL	ALE Inactive Delay (Note 8)			85		55	пş
TLLAX	Address Hold Time to ALE Inactive (Note 7)		59		34		ns
TCLDV	Data Valid Delay (Note 8)		10	110	10	60	ns
TCHDX	Data Hold Time (Note 10)		10		10		ns
TWHDX	Data Hold Time After WR (Note 9)		88		38		ns
TCVCTV	Control Active Delay 1 (Note 8)		10	110	10	70	ns
TCHCTV	Control Active Delay 2 (Note 8)	C _L = 100 pF	10	110	10	60	ns
TCVCTX	Control Inactive Delay (Note 8)	for all 8086 Outputs (in addition	10	110	10	70	ns
TAZRL	Address Float to READ Active (Note 9)	to 8086 internal loads)	0		0		ns
TCLRL	RD Active Delay (Note 8)		10	165	10	100	ns
TCLRH	RD Inactive Delay (Note 8)		10	150	10	80	ns
TRHAV	RD Inactive to Next Address Active (Note 10)		155		85		ns
TCLHAV	HLDA Valid Delay (Note 8)		10	160	10	100	ns
TRLRH	RD Width (Note 10)		325		200		ns
TWLWH	WR Width (Note 10)		340		210		ns
TAVAL	Address Valid to ALE LOW (Note 9)		58		28	Ι	ns
TOLOH	Output Rise Time (Note 9)	From 0.8 to 2.0 V		20		20	ns
TOHOL	Output Fall Time (Note 9)	From 2.0 to 0.8 V		12		12	ns
2. 3. 4. 5. 6. 7. 8. 9.	Signal at 8284A and 8288 shown for reference only. Setup requirement for asynchronous signal only to guara Applies only to T3 and wait states. Applies only to T2 state (8 ns into T3). Not tested; these specs are controlled by the Teradyne VCC — 4.5 V V VIH = 24 V VILC — 4.5 V VIHC = 4.3 V VILC = 25 V VOH = 1.6 V VOL = 1.4 V VOL = 1.6 V VCM Minimum spec tested at VCC Min. (4.5 V) only. Tested at VCC Min. (4.5 V) only. Tested at VCC Min. (4.5 V) only. Tested conditions for TCLCL Max. are:		CLK.				



SWITCHING CHARACTERISTICS over MILITARY operating range (continued)

MAX MODE SYSTEM (USING 8288 BUS CONTROLLER) TIMING REQUIREMENTS

Parameter	Parameter	Test Conditions	80	86	80	86-2	1
Symbol	Description	(Note 6)	Min.	Max.	Min.	Max.	Unit
TCLCL	CLK Cycle Period (Note 11)		200	500	125	500	ns
TCLCH	CLK LOW Time		118		68		ns
TCHCL	CLK HIGH Time		69		44		ns
TCH1CH2	CLK Rise Time (Note 5)	From 1.0 to 3.5 V		10		10	ns
TCL2CL1	CLK Fall Time (Note 5)	From 3.5 to 1.0 V		10		10	пѕ
TDVCL	Data in Setup Time		30		20	 	ns
TCLDX	Data in Hold Time		10		10		ns
TR1VCL	RDY Setup Time into 8284A (Notes 1 & 2)		35		35		ns
TCLR1X	RDY Hold Time into 8284A (Notes 1 & 2)		0		0		ns
TRYHCH	READY Setup Time into 8086		118		68		ns
TCHRYX	READY Hold Time into 8086		30		20		ns
TRYLCL	READY Inactive to CLK (Note 4)		-8		-8		ns
TINVCH	Setup Time for Recognition (INTR, NMI, TEST) (Note 2)		30		15		ns
TGVCH	RQ/GT Setup Time		30		15		ns
TCHGX	RQ Hold Time into 8066		40		30		ns
TILIH	Input Rise Time (Except CLK) (Note 5)	From 0.8 to 2.0 V		20		20	ns
TIHIL	Input Fall Time (Except CLK) (Note 5)	From 2.0 to 0.8 V		12		12	ns

Notes: 1. Signal at 8284A and 8288 shown for reference only.

2. Setup requirement for asynchronous signal only to guarantee recognition at next CLK.

3. Applies only to T3 and wait states.

4. Applies only to T2 state (8 ns into T3).

5. Not tested; these specs are controlled by the Teradyne J941 tester.

6. Voc = 4.5 V VIH = 2.4 V
VIL = .45 V VIHC = 4.3 V
VILC = .25 V VOH = 1.6 V
VOL = 1.4 V

7. Minimum spec tested at Voc Max. (5.5 V) only.

8. Maximum spec tested at Voc Min. (4.5 V) only.

10. Tested at Voc Min. (4.5 V) only.

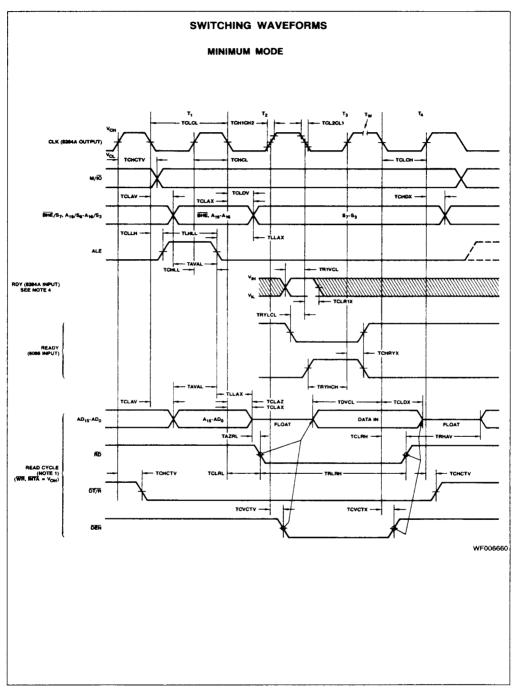
11. Test conditions for TCLCL Max. are:
Voc = 4.5 V VIH = 4 V
VILC = 0 V VIH = 4 V
VILC = 0 V VIHC = 5 V



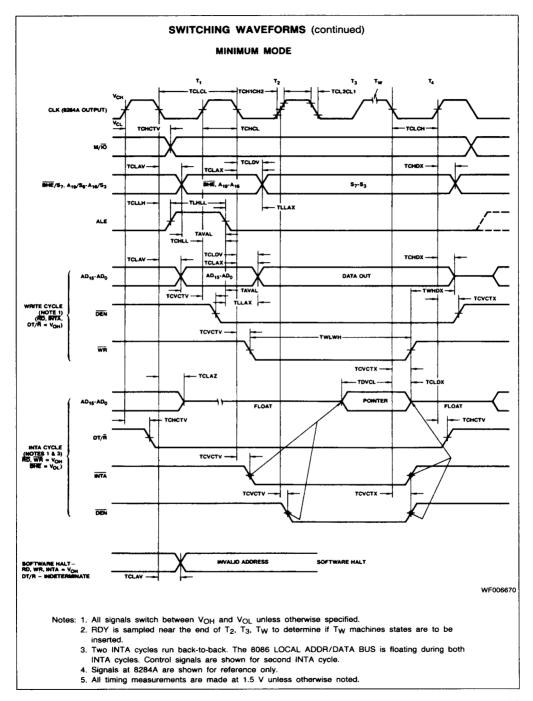
SWITCHING CHARACTERISTICS over MILITARY operating range (continued) TIMING RESPONSES

		Test	80	86	801	36-2	
Parameter Symbol	Parameter Description	Conditions (Note 6)	Min.	Max.	Min.	Max.	Unit
TCLML	Command Active Delay (Note 1)		10	35	10	35	ns
TCLMH	Command Inactive Delay (Note 1)		10	35	10	35	ns
TRYHSH	READY Active to Status Passive (Note 3)	1		110		65	ns
TCHSV	Status Active Delay (Notes 7 & 8)	1	10	110	10	60	ns
TCLSH	Status Inactive Delay		10	130	10	70	ns
TCLAV	Address Valid Delay	1	10	110	10	60	ns
TCLAX	Address Hold Time		10		10		ns
TCLAZ	Address Float Delay	1	10	80	10	50	ns
TSVLH	Status Valid to ALE HIGH (Note 1)	1		15		15	ns
TSVMCH	Status Valid to MCE HIGH (Note 1)			15		15	ns
TCLLH	CLK LOW to ALE Valid (Note 1)			15		15	ns
TCLMCH	CLK LOW to MCE HIGH (Note 1)			15		15	ns
TCHLL	ALE Inactive Delay (Note 1)	C _L = 100 pF for all 8086		15		15	ns
TCLMCL	MCE Inactive Delay (Note 1)	Outputs (In addition to 8086 internal loads)		15		15	ns
TCLDV	Data Valid Delay		10	110	10	60	ns
TCHDX	Data Hold Time		10		10		ns
TCVNV	Control Active Delay (Note 1)		5	45	5	45	ns
TCVNX	Control Inactive Delay (Note 1)		10	45	10	45	ns
TAZRL	Address Float to Read Active		0		0		ns
TCLRL	RD Active Delay		10	165	10	100	ns
TCLAH	RD Inactive Delay]	10	150	10	80	ns
TRHAV	RD Inactive to Next Address Active		155		85		ns
TCHDTL	Direction Control Active Delay (Note 1)			50		50	ns
TCHDTH	Direction Control Inactive Delay (Note 1)]		30		30	ns
TCLGL	GT Active Delay (Note 8)]	0	85	0	50	nş
TCLGH	GT Inactive Delay (Note 8)	1	0	85	0	50	ns
TRLRH	RD Width	1	325		200		ns
TOLOH	Output Rise Time	From 0.8 to 2.0 V		20		20	ns
TOHOL	Output Fall Time	From 2.0 to 0.8 V		12		12	ns

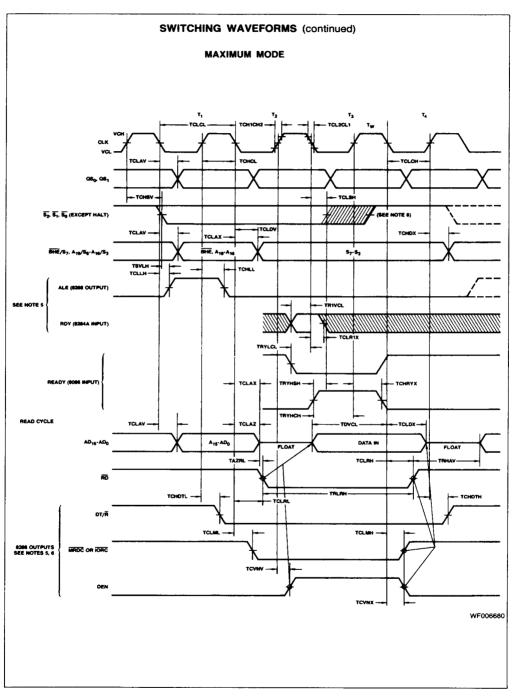
Notes: 1. Signal at 8284A and 8288 shown for reference only.
2. Setup requirement for asynchronous signal only to guarantee recognition at next CLK.
3. Applies only to T3 and wait states.
4. Applies only to T2 state (8 ns into T3).
5. Not tested; these specs are controlled by the Teradyne J941 tester.
6. Vcc = 4.5 V, 5.5 V VIH = 2.4 V
VIL = .45 V VIHC = 4.3 V
VILC = .25 V VOH = 1.6 V
VOL = 1.4 V
7. Minimum spec tested at Vcc Max. (5.5 V) only.
8. Maximum spec tested at Vcc Min. (4.5 V) only.
9. Tested at Vcc Min. (4.5 V) only.
10. Tested at Vcc Min. (4.5 V) only.
11. Test conditions for TCLCL Max. are:
Vcc = 4.5 V
VIL = 0 V VIH = 4 V
VILC = 0 V VIHC = 5 V



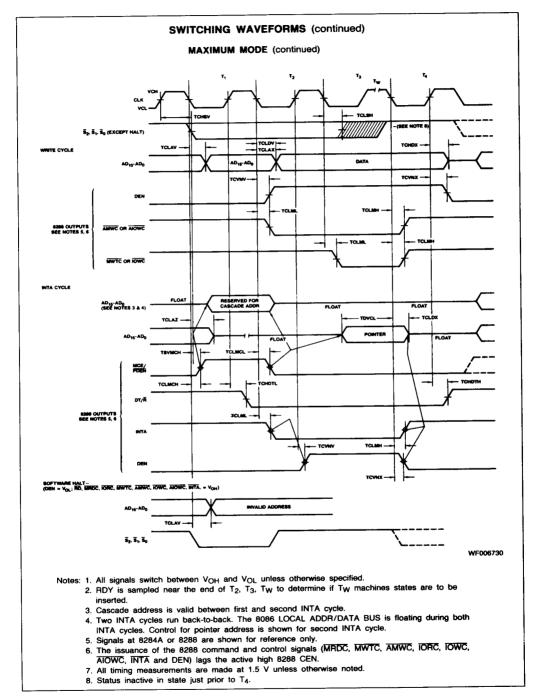
1-24 8086



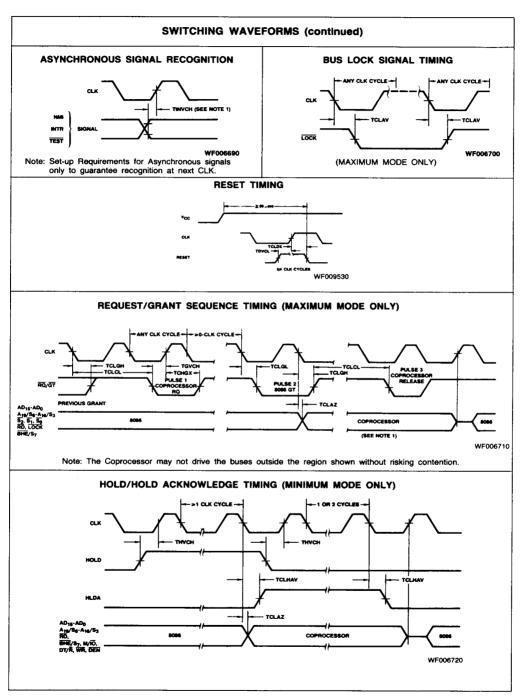
8086



1-26 8086







1-28 8086

	8086/808 INSTRUCTION SET	8 SUMMARY		
DATA TRANSFER				
MOV = Move	76543210	76543210	7 6 5 4 3 2 1 0	7654321
Register/memory to /from register	100010dw	mod reg r/m		
mmediate to register/memory	1100011w	mod 0 0 0 r/m	data	data if w = 1
Immediate to register	1 0 1 1 w reg	data	data if w = 1	
Memory to accumulator	1010000w	addr-low	addr-high]
Accumulator to memory	1010001w	addr-low	addr-high]
Register/memory to segment register	10001110	mod 0 reg r/m		
Segment register to register/memory	10001100	mod 0 reg r/m		
PUSH = Push:				
Register/memory	1111111	mod 1 1 0 r/m		
Register	0 1 0 1 0 reg			
Segment register	0 0 0 reg 1 1 0			
POP = Pop:				
Register/memory	10001111	mod 0 0 0 r/m		
Register	0 1 0 1 1 reg			
Segment register	0 0 0 reg 1 1 1			
XCHG = Exchange:	T		ı	
Register/memory with register	1000011w	mod reg r/m		
Register with accumulator	1 0 0 1 0 reg			
IN = Input from:				
Fixed port	1110010w	port		
Variable port	1110110 w			
OUT - Ouput to:			•	
Fixed port	1110011 w	port	ļ	
Variable port	1110111W			
XLAT - Transtate byte to AL	11010111		1	
LEA = Load EA to register	10001101	mod reg r/m	ļ	
LDS = Load pointer to DS	11000101	mod reg r/m]	
LES - Load pointer to ES	11000100	mod reg r/m	J	
LANF - Load AH with flags	10011111			
SANF - Store AH into flags	10011110			
PUSHF = Push flags	10011100			
POPF = Pop flags	10011101			



	TRUCTION SET SUM	MINIART (CONTIN	iuea)	
ARITHMETIC				
ADD = Add	7 6 5 4 3 2 1 0	76543210	76543210	76543210
Reg/memory with register to either	00000dw	mod reg r/m		
Immediate to register / memory	100000sw	mod 0 0 0 r/m	data	data if s:w = 01
Immediate to accumulator	0000010w	data	data if w = 1	
ADC = Add with carry:				
Reg/memory with register to either	000100dw	mod reg r/m		
Immediate to register/memory	100000sw	mod 0 1 0 r/m	data	data if s:w = 01
mmediate to accumulator	0 0 0 1 0 1 0 w	data	data if w = 1	
INC = Increment:				
Register/memory	111111w	mod 0 0 0 r/m		
Register	0 1 0 0 0 reg			
AAA - ASCII adjust for add	00110111			
DAA = Decimal adjust for add	00100111			
SUB = Subtract:				
Reg/memory and register to either	001010dw	mod reg r/m		
mmediate from register/memory	100000sw	mod 1 0 1 r/m	data	data if s:w = 01
mmediate from accumulator	0 0 1 0 1 1 0 w	data	data if w = 1	
SBB = Subtract with borrow:				
Reg/memory and register to either	000110dw	mod reg r/m		
mmediate from register/memory	100000sw	mod 0 1 1 r/m	data	data if s:w = 01
mmediate from accumulator	0001110w	data	data if w = 1	uuu
DEC = Decrement:				
Register/memory	111111W	mod 0 0 1 r/m		
Register	0 1 0 0 1 reg			
ŧEG Change sign	1111011w	mod 0 1 1 r/m		
CMP = Compare:	[0011101]			
Register/memory with register	0011101w	mod reg r/m		
Register with register/memory	0011100w	mod reg r/m		
mmediate with register/memory	100000sw	mod 1 1 1 r/m	data	data if s:w = 01
mmediate with accumulator	0011110w	data	data if w = 1	
AS ASCII adjust for subtract	0011111			
DAS Decimal adjust for subtract	00101111			
IUL Mulitiply (unsigned)	1111011w	mod 1 0 0 r/m		
MUL Integer multiply (signed):	1111011w	mod 1 0 1 r/m		
AM ASCII adjust for multiply	11010100	00001010		
IV Divide (unsigned):	1111011w	mod 1 1 0 r/m		
DIV Integer divide (signed)	1111011w	mod 1 1 1 r/m		
AD ASCH adjust for divide	11010101	00001010		
SW Convert byte to word	10011000			
WD Convert word to double word	10011001			

INSTRUCTION SET SUMMARY (continued) LOGIC 76543210 76543210 76543210 76543210 1 1 1 1 0 1 1 w | mod 0 1 0 r/m NOT Invert 1 1 0 1 0 0 v w | mod 1 0 0 r/m SHL/SAL Shift logical/arithmetic left 1 1 0 1 0 0 v w | mod 1 0 1 r/m SHR Shift logical right 1 1 0 1 0 0 v w | mod 1 1 1 r/m SAR Shift arithmetic right 1 1 0 1 0 0 v w | mod 0 0 0 r/m ROL Rotate left 110100vw mod 0 0 1 r/m ROR Rotate right 110100vw mod 0 1 0 r/m RCL Rotate through carry flag left mod 0 1 1 r/m 110100vw RCR Rotate through carry right AND = And: 001000dw mod reg r/m Reg/memory and register to either data if w = 1 1000000w mod 1 0 0 r/m data Immediate to register/memory data if w = 1 0010010w data Immediate to accumulator TEST = And function to flags, no result: 1000010w mod reg r/m Register/memory and register data if w = 1 mod 0 0 0 r/m data 1111011w Immediate data and register/memory data if w = 1 1010100w data Immediate data and accumulator OR = Or: 000010dw mod reg r/m Reg/memory and register to either data if w = 1 mod 0 0 1 r/m data 1000000w Immediate to register/memory data if w = 1 0000110w data Immediate to accumulator XOR = Exclusive or: 001100dw mod reg r/m Reg/memory and register to either data if w = 1 mod 1 1 0 r/m data 1000000w Immediate to register/mømory data if w = 1 0011010w data Immediate to accumulator STRING MANIPULATION: 1111001z REP - Repeat 1010010w MOVS - Move byte/word 1010011w CMPS = Compare byte/word 1010111w SCAS = Scan byte/word 1010110w LODS - Load byte/wd to AL/AX 1010101w STOS = Stor byte/wd from AL/A



CALL = Call	76543210	76543210	76543210	76543216
Direct within segment	11101000	disp-low	disp-high]
indirect within segment	1111111	mod 0 1 0 r/m		•
Direct intersegment	10011010	offset-low	offset-high]
		seg-low	seg-high	
Indirect intersegment	11111111	mod 0 1 1 r/m		•
JMP = Unconditional jump:				
Direct within segment	11101001	disp-low	disp-high	1
Direct within segment-short	11101011	disp		
Indirect within segment	1111111	mod 1 0 0 r/m		
Direct intersegment	11101010	offset-low	offset-high	1
		seg-low	seg-high	
Indirect intersegment	1111111	mod 1 0 1 r/m		
RET = Return from CALL:				
Within segment	11000011			
Within seg adding immed to SP	11000010	data-low	data-high	
ntersegment	11001011		outu riigii	
ntersegment adding immediate to SP	11001010	data-low	data-high	
E/JZ = Jump on equal/zero	01110100	disp		
IL/JNGE = Jump on less/not greater or equal	01111100	disp		
JLE/JNG = Jump on less or equal/not greater	0111110	disp		
JB/JNAE = Jump on below/not above or equal	01110010	disp		
BE/JNA = Jump on below or equal/not above	01110110	disp		
IP/JPE = Jump on parity/parity even	01111010	disp		
O = Jump on overflow	01110000	disp		
S = Jump on sign	01111000	disp		
INE/JNZ = Jump on not equal/not zero	01110101	disp		
INL/JGE = Jump on not less/greater or equal	01111101	disp		
NLE/JG = Jump on not less or equal/greater	0111111	disp		
NB/JAE = Jump on not below/above or equal	01110011	disp		
NBE/JA = Jump on not below or equal/above	01110111	disp		
NP/JPO = Jump on not par/par odd	01111011	disp		
NO - Jump on not overflow	01110001	disp		
NS = Jump on not sign	01111001	disp		
OOP = Loop CX times	11100010	disp		
OCPZ/LOOPE = Loop while zero/equal	11100001	disp		
OOPNZ/LOOPNE = Loop while not zero/equal	11100000	disp		
CXZ - Jump on CX zero	11100011	disp		

INSTRUCTION SET SUMMARY (continued) CONTROL TRANSFER (Cont'd.) 76543210 76543210 76543210 76543210 INT = Interrupt 11001101 type Type specified 11001100 Type 3 11001110 INTO = Interrupt on overflow 11001111 IRET = interrupt return PROCESSOR CONTROL 11111000 CLC = Clear carry 11110101 CMC = Complement carry 11111001 STC = Set carry 11111100 CI D = Clear direction 11111101 STD = Set direction 11111010 CLL = Clear interrupt 11111011 ST! = Set interrupt 11110100 HLT - Halt 10011011 WAIT - Wait 11011xxx mod x x x r/m ESC = Processor Extension Escape 11110000 LOCK = Bus lock prefix

Footnotes: AL = 8-bit accumulator AX = 16-bit accumulator

AX = 16-bit accumulator
CX = Count register
DS = Data segment
ES = Extra segment
ES = Extra segment
Above/below refers to unsigned value.
Greater = more positive.
Less = less positive (more negative) signed values
if d = 1 then "to" reg; if d = 0 then "from" reg
w = 1 then word instruction; if w = 0 then byte instruction

if mod = 11 then r/m is treated as a REG field
if mod = 00 then DISP = 0 , disp-low and disp-high are absent
if mod = 01 then DISP = disp-low sign-extended to 16-bits, disp-high is
absent
if mod = 10 then DISP = disp-low sign-extended to 16-bits, disp-high is
if r/m = 000 then EA = (BX) + (SI) + DISP
if r/m = 010 then EA = (BP) + (SI) + DISP
if r/m = 100 then EA = (BP) + (SI) + DISP
if r/m = 101 then EA = (BP) + (DI) + DISP
if r/m = 101 then EA = (BP) + (DI) + DISP
if r/m = 101 then EA = (BP) + (DI) + DISP
if r/m = 101 then EA = (BP) + (DI) + DISP
if r/m = 101 then EA = (BP) + (DI) + DISP
if r/m = 111 then EA = (BP) + DISP
if r/m = 111 then EA = (BY) + DISP

DISP follows 2nd byte of instruction (before data if required)
*except if mod = 00 and r/m = 110 then EA = disp-high: disp-low.

if s:w = 01 then 16 bits of immediate data form the operand.
if s:w = 11 then an immediate data byte is sign extended to form the
16-bit operand.

if v = 0 then "count" = 1; if v = 1 then "count" in (CL) x = don't care

z is used for string primitives for comparison with ZF Flag.

SEGMENT OVERRIDE PREFIX

0 0 1 reg 1 1 0							
	0	0	1	reg	1	1	0

REG is assigned according to the following table:

16-Bit (w = 1)	8-Bit (w = 0)	Segment
000 AX	000 AL	00 ES
001 CX	001 CL	01 CS
010 DX	010 DL	10 SS
011 BX	011 BL	11 DS
100 SP	100 AH	
101 BP	101 CH	
110 SI	110 DH	
111 DI	111 BH	

Instructions which reference the flag register files as a 16-bit object use the symbol FLAGS to represent the file:

FLAGS = X:X:X:X:(OF):(DF):(TF):(SF):(ZF):X:(AF):X:(PF):X:(CF)